## COMPUTER SCIENCE 322 (Winter Term 2004) Compiler Construction

Prof. Levy

## Problem Set 6

Due Monday 08 March

Reading Assignment: Dragon Book 4.7; 7.6

## Written Assignment

Part I: Dragon Book exercise 4.33 (e), (f)

Part II: Recall our grammar for let expressions in ML, reproduced below with additional rules for arithmetic and boolean expressions:<sup>1</sup>

```
expr 	o let-expr \mid arith-expr \mid bool-expr \mid (expr) \mid id \mid number \mid boolean \ let-expr 	o let \ bindings \ in \ expr \ end \ arith-expr 	o expr + expr \mid expr * expr \ bool-expr 	o expr \ and \ expr \mid expr \ or \ expr \ bindings 	o binding \mid binding; \ bindings \ binding 	o val \ var = expr
```

with the usual assumptions about precedence (\* > +, and > or), and the following regular definitions:

```
id \rightarrow [A-Za-z] [A-Za-z0-9]*
number \rightarrow [0-9]+
boolean \rightarrow true | false
```

Suppose someone makes the following suggestion: for this sort of simple programming language, we can avoid writing a distinct type-checking component, by incorporating type information into the grammar itself. For example, to prohibit mixing up booleans and numbers, we could move **number** from the RHS of the *expr* rule to the RHS of the *arith-expr*, and move **boolean** to the RHS of the *bool-expr* rule.

What is wrong with this argument? Specifically, what crucial element does the solution fail to handle?

Turn your answers in to me on paper.

<sup>&</sup>lt;sup>1</sup>in reality, ML uses **andalso** instead of **and**, and **orelse** instead of **or**, but I like the simpler versions better.

## Programming Assignment

This assignment has two purposes: (1) familiarizing yourself with building parsers using the JFlex/CUP alternative to lex/yacc; (2) trying out different symbol table implementations to see which are most efficient. To complete the assignment, perform the following steps:

- 1. Unzip the file ps6.zip
- 2. Write a CUP grammar called TinyML.cup, based on the grammar in the written assignment. Your grammar should include the precedences mentioned above.
- 3. Complete the lexical rules for the JFlex scanner in TinyML.flex, using the appropriate symbol methods for the return values (all three can have hollow methods to start):
- 4. Following the SymbolTable interface, implement three different versions of a symbol table:
  - (a) LinearSymbolTable, which uses a simple list (e.g., java.util.Vector) to store the symbol table entries.
  - (b) PJWHashedSymbolTable, which uses a hashtable based on the PJW hash-function algorithm in Fig. 7.35 on page 4.26 of the Dragon book. *Hint*: use long instead of unsigned (which Java doesn't have), to avoid numerical overflow.
  - (c) SunHashedSymbolTable, which uses Sun's java.util.Hashtable (one-liner's for each method you implement).
- 5. Use the gentest program provided to generate sample inputs for your parser. Running the command

```
\% java gentest N>\mathit{outfile}
```

will put a little ML program with N randomly-named variable declarations into the file outfile. You can then test your parser on this file by running the command

% java TinyML outfile. Compare the time taken by the three symbol table implementations, using enough different values of N.

- 6. Turn in the following items:
  - (a) TinyML.cup
  - (b) TinyML.flex
  - (c) LinearSymbolTable.java
  - (d) PJWHashedSymbolTable.java
  - (e) SunHashedSymbolTable.java
  - (f) Either on paper or in a file, a graph or table showing your results from the last step.